Exploiting Wireless Channel State Information for Throughput Maximization

V. Tsibonis, L. Georgiadis, L. Tassiulas

Abstract— **We consider the problem of scheduling packets over a number of channels with time varying connectivity. Policies proposed for this problem either stabilize the system when the arrival rates are within the stability region, or optimize an objective function under the assumption that all channel queues are saturated. We address the realistic situation where it is not known apriori whether the channel queues are saturated or not, and provide a scheduling policy that maximizes the weighted sum of channel throughputs. We employ a burstiness-constrained channel model that allows us to dispense of statistical assumptions and simplifies the proofs.**

Index Terms— **Scheduling, Deterministic network calculus, QoS in wireless networks.**

I. INTRODUCTION

The primary motivation of this work is the problem of scheduling the transmissions of multiple data flows sharing the same wireless channel. The relative delay tolerance of data applications, together with the bursty traffic characteristics, opens up the potential for scheduling transmissions so as to optimize throughput. Given the above considerations, we examine a time-slotted parallel queue system with a single server. The condition of the associated channel of every queue varies with time between "on" and "off" states. In every time slot only one packet can be transmitted from a given queue, if the associated channel is in the "on" state and the queue is non empty. The main result of this paper is the design of a scheduling policy that allocates the server to the queues in such a way that the weighted sum of channel throughputs is maximal.

A related approach along these lines is proposed in [4], where the authors identify optimality properties for scheduling downlink transmissions to data users, in CDMA networks. For arbitrary-topology networks, the problem of admission control and rate allocation to the users so that certain quality-of-service requirements are met, is investigated. A mathematical programming formulation is obtained for determining the optimal transmission schedule. The effect of wireless channels on the performance of transmission protocols such as TCP is examined through simulations in [5]. The authors conclude that channel-state scheduling can lead to significant improvement in channel utilization.

The problem of scheduling wireless channels with time varying connectivity has been addressed in the past in several

V. Tsibonis and L. Georgiadis are with Aristotle University of Thessaloniki, Thessaloniki, GREECE. e-mails: vtsib@egnatia.ee.auth.gr, leonid@eng.auth.gr

L. Tassiulas is with the University of Thessaly, Greece and University of Maryland, College Park. e-mail: leandros@Glue.umd.edu

different contexts. In [12], optimal scheduling for a wireless system consisting of multiple queues and a single server is studied. The arrival processes to the queues are assumed i.i.d Bernoulli. The wireless channels can be in the "on" or in the "off" state according to i.i.d Bernoulli processes. The authors derived the system stability region; moreover, they showed that the policy that among the queues whose channel is "on" serves the one with the longest queue, stabilizes the system whenever the arrival rates are within the stability region. In [13], Tassiulas considered a system that generalizes the one in [12] in the following aspects. First, a network with arbitrary topology is considered. Second, the topology is represented by a hidden Markov model instead of an independent and identically distributed (i.i.d) process. Third, anticipative scheduling policies are taken into consideration. Fourth, multiple link transmission rates are considered. In that context, after the characterization of the region of achievable throughputs, a transmission scheduling policy is proposed, that achieves all throughput vectors achievable by any anticipative policy.

The problem of scheduling transmissions over a wireless channel with time-varying transmission rates is considered in [11], [6], [9] and [8]. The problem of providing a scheduling policy that stabilizes the system whenever the arrival-rate vector lies within the stability region is dealt in [11] and in [8]. In [11], a finite set of channel states is assumed and every channel can be in one of these states. With each state there is an associated data rate, representing the rate at which the queue is served if selected for transmission. The arrival processes to the queues are assumed mutually independent, ergodic, Markov chains with countable state space. Under these assumptions it was shown that the scheduling policy, called the exponential rule, makes the queues stable if there exists any policy that can do so. In [8], the authors consider the problem of power and server allocation in a multi-beam satellite downlink which transmits data to different ground locations over time varying channels. The authors establish the stability region of the system and develop a power allocation policy, which stabilizes the system whenever the system is stabilizable and when the arrival and channel state processes are i.i.d.

In [9] and [6] the problem of developing a scheduling policy for efficient channel utilization is considered for the case that all the queues are unstable. In [9] the state of a channel is modeled by a stochastic process, which represents the level of performance of the given channel. A scheduling policy is provided which maximizes the average system performance given that a predetermined time-fraction assignment

is achieved when all the queues are unstable. In [6], the authors consider a base station serving data-users. The feasible rates of the users vary over time according to some stationary discrete-time stochastic process. A scheduling policy that exploits the variations in the channel conditions and maximizes the minimum throughput is developed.

The main contribution of this paper is that it studies a system with time-varying connectivity when the arrival rates do not belong necessarily in the stability region of the system. This is an important situation that arise in practice, since the channel parameters and the arrival rates may not be known apriori, or may vary over time. In such a case, scheduling policies proposed before for maximizing the throughput may fail, and the system may have a rather erratic behavior. In the current work we consider the scheduling problem of maximizing the weighted sum of user throughputs. We provide a scheduling policy that is optimal under any arrival rates. In the most general case, under the optimal policy we propose, some queues will be stable while others will operate in saturation. Such a dynamic behavior makes the analysis of the system rather difficult. Instrumental in the analysis of our policy was the adoption of a "bounded burstiness" model for the variability of the channel inspired by "burstiness constrained" traffic models that have been used over the last several years in the analysis of rate-controlled communication networks [2].

The paper is organized as follows. In Section II, the traffic and channel model is introduced. Specifically, the constraints on the arrival and slot availability processes are given. In Section III we provide the problem formulation and define the scheduling policy. In Section IV we provide the optimality proof of the proposed policy. Conclusions and suggestions for further work are discussed in Section V.

A. Notations and Conventions

Before we proceed, we discuss some of the notations and conventions that we use throughout the paper. Sets of numbers are denoted by calligraphic capital letters. In particular we define $\mathcal{N} = \{1, ..., N\}$. A subset $\mathcal{S}$ of a set $\mathcal{D}$ is denoted by $\mathcal{S} \subseteq \mathcal{D}$ and a strict subset by $\mathcal{S} \subset \mathcal{D}$. In several places we will use sets as subscripts or arguments, say $F(\mathcal{S})$. To simplify notation and if there is no possibility for confusion, instead of $F(\{i_1, ..., i_k\})$ we write $F(i_1, ..., i_k)$. Also, we write $\sum_{\mathcal{S}} y_i$ to denote $\sum_{i \in \mathcal{S}} y_i$. If $\mathcal{S} = \varnothing$, then we define $\sum_{\mathcal{S}} x_i = 0$. Also, $\cup_{i=k}^{l} \mathcal{D}_i = \varnothing$ if $k > l$. The cardinality of a set $\mathcal{S}$ is denoted by $|\mathcal{S}|$. If $\mathbf{X} = [x_{ij}]$ and $\mathbf{Y} = [y_{ij}]$ are matrices, then $\mathbf{X} \leq \mathbf{Y}$ ($\mathbf{X} < \mathbf{Y}$) means that $x_{ij} \leq y_{ij}$ ($x_{ij} < y_{ij}$) for all i and j. Finally, by $\mathbf{X}^T$ we denote the transpose of $\mathbf{X}$.

Some of the rather technical proofs in Sections III and IV that are not essential for understanding the main ideas of the arguments, are omitted due to lack of space. These proofs can be found at the site http://genesis.ee.auth.gr /SITE_AUTH_UNIVERSITY/SITE_TDIVISION /users/georgiadis/english/PersonalPages /ConferencePapers/TimeVarInfoc03.pdf.

II. TRAFFIC AND CHANNEL MODEL

We consider a system consisting of N channels. With each channel there is an associated queue holding packets that are to be transmitted over the given channel. Packets are of fixed size and time is divided in slots of unit length, equal to the transmission time of a packet. Slot $t \geq 1$ refers to the interval $(t-1, t]$. In the interval $(t-1, t]$ (slot t), $a_i(t)$ new packets join queue i to be transmitted over the corresponding channel. At the beginning of slot t, i.e., at time $t - 1$, one packet among those already in one of the N queues may be chosen for transmission at slot t. The number of packets from queue i transmitted in slot t is $b_i(t)$ (therefore, $b_i(t)$ is either 0 or 1) and the number of packets in queue i at time $t \geq 0$ is $q_i(t)$. Therefore the number of packets at queue i, $i \in \mathcal{N}$, evolves with time according to the equation

$$q_i(t) = (q_i(t-1) - b_i(t))^+ + a_i(t),$$

where $(x)^+ = \max\{0, x\}$.

Define $a_{\mathcal{S}}(t) = \sum_{\mathcal{S}} a_i(t)$ and $b_{\mathcal{S}}(t) = \sum_{\mathcal{S}} b_i(t)$.

At slot t, channel i may or may not be available for transmission of queue i packets. If the channel is available for transmission, we say that the channel is in the "on" state. We define for $\mathcal{S} \subseteq \mathcal{N}$, $\mathcal{S} \neq \varnothing$,

$$c_{\mathcal{S}}(t) = \begin{cases} 1 & \text{if at least one channel in } \mathcal{S} \text{ is} \\ & \text{on in slot } t \\ 0 & \text{otherwise} \end{cases}$$

and $c_{\varnothing}(t) \equiv 0$. For example, Figure 1 shows the channel availability for 3 channels during 15 time slots. According to the figure

- $c_{\{1\}}(t) = 1$ for $1 \leq t \leq 12$ and zero elsewhere.
- $c_{\{3\}}(t) = 1$ for $2 \leq t \leq 7$, $9 \leq t \leq 14$ and zero elsewhere.
- $c_{\{1,2\}}(t) = 1$ for $1 \leq t \leq 15$ and zero elsewhere.
- $c_{\{1,3\}}(t) = 1$ for $1 \leq t \leq 14$ and zero elsewhere.
- $c_{\{1,2,3\}}(t) = 1$ for $1 \leq t \leq 15$ and zero elsewhere.

Since only one packet may be transmitted in one slot, we have

$$b_{\mathcal{S}}(t) = \begin{cases} 1 & \text{if } b_i(t) = 1 \text{ for one of the} \\ & \text{channels in } \mathcal{S} \\ 0 & \text{otherwise} \end{cases}.$$

Transmission over channel i make take place ($b_i(t) = 1$) only if the channel is in the "on" state and hence,

$$b_{\mathcal{S}}(t) \leq c_{\mathcal{S}}(t). \tag{1}$$

If $x(t)$ is any of the quantities defined above, we denote

$$X(s, t) = \sum_{\tau=s+1}^{t} x(\tau).$$

We make the following assumptions regarding the traffic and slot availability processes.

Traffic Model:

$A_i(s, t)$ is $(\sigma_i^U, \sigma_i^L, \alpha_i)$-constrained, i.e., it holds for any $t \geq s \geq 0$,

$$\alpha_i(t - s) - \sigma_i^L \leq A_i(s, t) \leq \alpha_i(t - s) + \sigma_i^U, \tag{2}$$

where

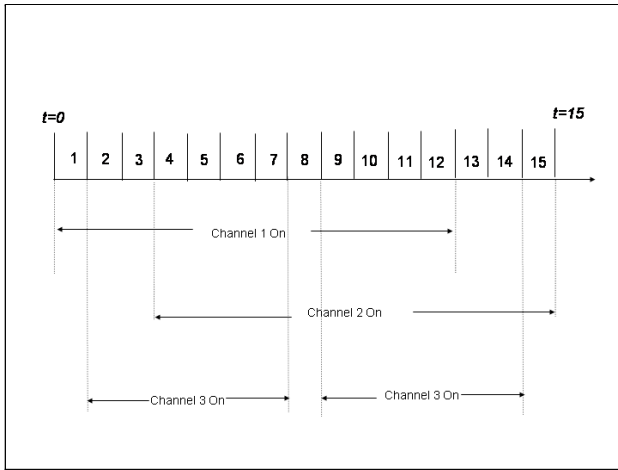$$\sigma_i^L \geq 0, \ \sigma_i^U \geq 0, \ \infty \geq \alpha_i \geq 0.$$

Fig. 1. Channel availability ("on" state) for 3 channels

Parameter α_i is the packet arrival rate to queue i (i.e., $\alpha_i = \lim_{t \to \infty} \frac{A_i(0,t)}{t}$) for transmission over the corresponding channel. We allow the possibility that $\alpha_i = \infty$, in order to include the case that some of the queues are infinite for $t \geq 1$. Note that it follows from the definition that if $A_i(s,t)$ is $(\sigma_i^U, \sigma_i^L, \alpha_i)$-constrained for $i \in \mathcal{S}$, then $A_{\mathcal{S}}(s,t)$ is $\left(\sum_{\mathcal{S}} \sigma_i^U, \sum_{\mathcal{S}} \sigma_i^L, \sum_{\mathcal{S}} \alpha_i\right)$-constrained.

Channel Availability Model:

$C_{\mathcal{S}}(s,t)$, is $(\theta_{\mathcal{S}}^U, \theta_{\mathcal{S}}^L, F(\mathcal{S}))$-constrained, i.e., it holds for any $t \geq s \geq 0$,

$$F(\mathcal{S})(t-s) - \theta_{\mathcal{S}}^L \leq C_{\mathcal{S}}(s,t) \leq F(\mathcal{S})(t-s) + \theta_{\mathcal{S}}^U, \tag{3}$$

where

$$\theta_{\mathcal{S}}^L \geq 0, \ \theta_{\mathcal{S}}^U \geq 0.$$

We also use the convention $F(\varnothing) = \theta_{\varnothing}^L = \theta_{\varnothing}^U = 0$.

We refer to the inequalities in (2) and (3) as "burstiness constraints". With the exception of α_i, $i \in \mathcal{N}$, the parameters used in the previous models are assumed finite.

The definitions for the traffic model are standard, see e.g., [2], [3], [1]. We elaborate on the Slot Availability Model. From (3) it follows that

$$\lim_{t \to \infty} \frac{C_{\mathcal{S}}(0,t)}{t} = F(\mathcal{S}), \tag{4}$$

that is, $F(\mathcal{S})$ is equal to the long-term fraction of time that at least one of the channels in $\mathcal{S}$ is in the "on" state. Also, from the definition of $c_{\mathcal{S}}(t)$ we have that for any subsets $\mathcal{S}, \mathcal{T}$, of $\mathcal{N}$, and for every t, it holds,

$$c_{\mathcal{T}}(t) \leq c_{\mathcal{S}}(t), \ \text{if } \mathcal{T} \subseteq \mathcal{S}$$
$$c_{\mathcal{S}}(t) + c_{\mathcal{T}}(t) \geq c_{\mathcal{S} \cup \mathcal{T}}(t) + c_{\mathcal{S} \cap \mathcal{T}}(t),$$

and hence

$$C_{\mathcal{T}}(s,t) \leq C_{\mathcal{S}}(s,t), \ \text{if } \mathcal{T} \subseteq \mathcal{S} \tag{5}$$
$$C_{\mathcal{S}}(s,t) + C_{\mathcal{T}}(s,t) \geq C_{\mathcal{S} \cup \mathcal{T}}(s,t) + C_{\mathcal{S} \cap \mathcal{T}}(s,t). \tag{6}$$

From (4), (5), (6) we conclude that $F(S)$ satisfies the following relations for any subsets $\mathcal{T}$ and $\mathcal{S}$ of $\mathcal{N}$.

$$F(\varnothing) = 0, \tag{7a}$$
$$F(\mathcal{T}) \leq F(\mathcal{S}), \ \mathcal{T} \subseteq \mathcal{S}, \tag{7b}$$
$$F(\mathcal{T}) + F(\mathcal{S}) \geq F(\mathcal{T} \cup \mathcal{S}) + F(\mathcal{S} \cap \mathcal{T}). \tag{7c}$$

The last property is known as the *submodularity* property.

As an example, suppose that the channel availability pattern in Fig. 1, is repeated indefinitely, i.e., we have a periodic channel availability process. Consider the first channel, i.e., $\mathcal{S} = \{1\}$. It holds

$$c_{\{1\}}(t) = \begin{cases} 1, & \text{for } 1 \leq t \leq 12 \\ 0, & \text{for } 13 \leq t \leq 15 \end{cases}$$

and $c_{\{1\}}(t + 15) = c_{\{1\}}(t)$, for every time-slot $t \geq 1$. Therefore we have

$$\left\lfloor \frac{t-s}{15} \right\rfloor 12 \leq C_{\{1\}}(s,t) \leq \left\lceil \frac{t-s}{15} \right\rceil 12, \ \text{or}$$
$$\frac{12}{15}(t-s) - 12 \leq C_{\{1\}}(s,t) \leq \frac{12}{15}(t-s) + 12.$$

In conjunction with definition (3), the above inequality states that $C_{\{1\}}(s,t)$ is $(\theta_{\{1\}}^U, \theta_{\{1\}}^L, F(1))$-constrained, with $\theta_{\{1\}}^U = \theta_{\{1\}}^L = 12$ and $F(1) = 12/15$, i.e., $F(1)$ is equal to the long-term fraction of time that the first channel is on. Similarly we have that $C_{\mathcal{S}}(s,t)$ is $(\theta_{\mathcal{S}}^U, \theta_{\mathcal{S}}^L, F(\mathcal{S}))$-constrained and according to the figure

- For $\mathcal{S} = \{3\}$, $\theta_{\{3\}}^U = \theta_{\{3\}}^L = 12$ and $F(3) = 12/15$.
- For $\mathcal{S} = \{1,2\}$, $\theta_{\{1,2\}}^U = \theta_{\{1,2\}}^L = 0$ and $F(1,2) = 1$.
- For $\mathcal{S} = \{1,3\}$, $\theta_{\{1,3\}}^U = \theta_{\{1,3\}}^L = 14$ and $F(1,3) = 14/15$.
- For $\mathcal{S} = \{1,2,3\}$, $\theta_{\{1,2,3\}}^U = \theta_{\{1,2,3\}}^L = 0$ and $F(1,2,3) = 1$.

We close this section with a few comments on the adopted traffic and channel models. The assumption that channels can be in two states only is made for technical reasons and we make no claims to practicality. This assumption is adopted in order to simplify the situation and get a better insight into the problem at hand. The adopted burstiness constrained models provide a clear description of system dynamics and make possible the analysis of the system using basically elementary (although not straightforward) techniques. Compared to introducing statistical assumptions for these models, there are both advantages and disadvantages. Note that the stationarity assumption is not needed in our model, although the existence of averages is implied. On the other hand deterministic rather than stochastic bounds on process fluctuations are imposed.

III. PROBLEM FORMULATION

Consider a scheduling policy π that at the beginning of slot t, i.e., at time $t-1$, decides which packet (if any) to transmit to one of the channels that are "on" in slot t. Let

$$r_i^\pi = \liminf_{t \to \infty} \frac{B_i^\pi(0,t)}{t},$$

be the "throughput" of channel i under policy π.

Given costs c_i, $i \in \mathcal{N}, c_1 \geq c_2 \geq .... \geq c_N \geq 0$, our objective is to determine a policy such that the weighted sum of throughputs

$$\sum_{\mathcal{N}} c_i r_i^{\pi},$$

is maximal.

Assume that the channel state at a given slot is known to the scheduler at the beginning of that slot and consider the following policy.

Scheduling Policy π^*

With queue i associate an index $I_i(q)$ of the form

$$I_i(q) = \begin{cases} q & \text{if } q \leq (N+1-i)T \\ (N+1-i)T & \text{if } q > (N+1-i)T \end{cases}$$

where $T > 0$. At time t, consider the nonempty queues whose channel is on. Among these queues, let i be the one with largest index $I_i(q_i(t))$ (if there are multiple such queues select one randomly). Transmit a packet from queue i at slot $t+1$.

Our objective is to show that for T large enough, policy π^* maximizes the weighted sum of throughputs, irrespective of whether the overall system is stable or not. Before we proceed, it is worth observing the following.

- Only the order of the costs c_i, $i \in \mathcal{N}$, not the actual values determine policy π^*. This situation is similar to the well-known μc-rule in queueing theory.
- The traffic and channel model parameters determine how large T should be chosen. In other words, the policy depends on these parameters only through T. Although estimates of T can be obtained through the analysis that follows, in practice the traffic and channel parameters may not be known beforehand. Of course, one can pick very large values of T but this implies larger delays and slower convergence. Hence, development of adaptive methods for determining T seems a more appropriate plausible way of choosing T.

A. Achievable Throughput Space and Related Linear Optimization Problem

Consider that the system operates under an arbitrary scheduling policy. From (1), the definitions of $B_{\mathcal{S}}(s,t)$, $C_{\mathcal{S}}(s,t)$ and (4) we have for any $\mathcal{S} \subseteq \mathcal{N}$,

$$\begin{aligned} F(\mathcal{S}) &= \lim_{t \to \infty} \frac{C_{\mathcal{S}}(0,t)}{t} \\ &\geq \liminf_{t \to \infty} \frac{B_{\mathcal{S}}^{\pi}(0,t)}{t} \\ &= \liminf_{t \to \infty} \frac{\sum_{\mathcal{S}} B_i^{\pi}(0,t)}{t} \\ &\geq \sum_{\mathcal{S}} \liminf_{t \to \infty} \frac{B_i^{\pi}(0,t)}{t} \\ &= \sum_{\mathcal{S}} r_i^{\pi}. \end{aligned} \tag{8}$$

In addition, the fact that $A_i(0,t) \geq B_i^{\pi}(0,t)$ and (2) imply that for any $i \in \mathcal{N}$, it holds

$$0 \leq r_i^{\pi} \leq \alpha_i. \tag{9}$$

From (8) and (9) we see that the maximum weighted sum of throughputs that can be achieved by any scheduling policy, cannot exceed the value of the following optimization problem:

Linear Optimization Problem.

$$\max_x \sum_{i=1}^{N} c_i x_i,$$

where for $\mathcal{N} = \{1, 2, ..., N\}$

$$\sum_{\mathcal{S}} x_i \leq F(\mathcal{S}), \ \mathcal{S} \subseteq \mathcal{N} \tag{10a}$$

$$x_i \leq \alpha_i, \ i = 1, ..., N \tag{10b}$$

$$x_i \geq 0, \ i = 1, ..., N \tag{10c}$$

and $F(\mathcal{S})$ satisfies (7a), (7b), (7c).

Let $\mathcal{N}_k = \{1, ..., k\}$, and $\mathcal{N}_0 = \varnothing$. It can be shown that the solution to the previous optimization problem is given recursively by

$$x_k^* = \min \left\{ \alpha_k, \min_{\mathcal{D} \subseteq \mathcal{N}_{k-1}} \left\{ F(k \cup \mathcal{D}) - \sum_{\mathcal{D}} x_i^* \right\} \right\}, \tag{11}$$

for $k = 1, ..., N$. The proof is omitted.

Our objective in the next section is to show that the scheduling policy π^* achieves the throughputs defined by (11) and therefore is optimal.

IV. OPTIMALITY PROOF

Since we deal only with policy π^* in this section, in order to simplify the notation we eliminate π^* from all related notations, e.g., we use r_i in place of $r_i^{\pi^*}$.

Before going into the details of the proof, we give the general idea of the approach. In the general case, it can be shown that under π^*, a subset $\mathcal{U}$ of the queues will grow to infinity, while the rest of the queues will receive the maximum possible throughput, i.e., we have $r_i = \alpha_i$, $i \in \mathcal{N} - \mathcal{U}$. The queues in $\mathcal{N} - \mathcal{U}$ are called "stable" queues, while those in $\mathcal{U}$ "unstable". It can be proved that for any stable queue i, we have $r_i = x_i^*$. To determine the throughputs of the unstable queues we first show that for T large enough, each of the stable queues fluctuates in a certain range around kT, $0 \leq k \leq N$. This fact and the manner the indices are used to determine the scheduling decisions, implies that $r_k = x_k^*$ for all unstable queues. We mention that in the course of the proof, the fact that $r_i = x_i^*$, is established by starting from the smallest indices and moving to the largest, rather than by first proving the result for the stable and then for the unstable queues.

The following lemma will be useful in the sequel.

Lemma 1: For any subsets $\mathcal{S}_1$, $\mathcal{S}_2$ of $\mathcal{N}$ and any t it holds

$$\sum_{\mathcal{S}_1} q_i(t) - \sum_{\mathcal{S}_2} q_i(t) \leq$$

$$\sum_{\mathcal{S}_1} q_i(t-1) - \sum_{\mathcal{S}_2} q_i(t-1) + \sum_{\mathcal{S}_1} \alpha_i + \sum_{\mathcal{S}_1} \sigma_i^U$$

$$+ F(\mathcal{S}_2) + \theta_{\mathcal{S}_2}^U.$$

Proof: This is immediate from the burstiness constraints on the arrival and slot availability processes. ∎

Following the definition in [12], we call a scheduling policy "Longest Connected Queue (LCQ) First", the scheduling policy which among the queues whose channels are "on", selects the one with the largest number of packets (if there are at least two such queues, pick one arbitrarily).

In the next lemmas we use the following notation. For $\mathcal{D} \subseteq \mathcal{N}$, denote by $q_{\mathcal{D}}^{(l)}(t)$ the lth maximum of $\{q_i(t)\}_{i \in \mathcal{D}}$ and by $\pi_l^{\mathcal{D}}(t)$ a permutation of the indices in $\mathcal{D}$ such that $q_{\pi_l^{\mathcal{D}}(t)}(t) = q_{\mathcal{D}}^{(l)}(t)$, $l \in \{1, ..., |\mathcal{D}|\}$. Hence $q_{\mathcal{D}}^{(1)}(t) = \max_{i \in \mathcal{D}}\{q_i(t)\}$ and $q_{\mathcal{D}}^{(|\mathcal{D}|)}(t) = \min_{i \in \mathcal{D}}\{q_i(t)\}$. Also, let $\mathcal{S}_{\mathcal{D}}^{(l)}(t) = \cup_{j=1}^{l} \pi_j^{\mathcal{D}}(t)$ (i.e., the set containing the l largest queues among the queues in $\mathcal{D}$) and $\overline{\mathcal{S}}_{\mathcal{D}}^{(l)}(t) = \mathcal{D} - \mathcal{S}_{\mathcal{D}}^{(l)}(t)$.

We use the term "a set $\mathcal{G}$ has priority at time t" when at time t, policy π^* chooses one of the packets of the queues in $\mathcal{G}$ (if any) for transmission, provided that the associated channel is on.

In order to avoid cluttering the notation in the proofs, in the following we will use the symbol $\mathbb{O}$, to denote a finite nonnegative quantity that depends only on the parameters of the arrival and slot availability processes, and $|\mathcal{N}|$. As will be clear from the proofs, in principle $\mathbb{O}$ can be explicitly computed - e.g., in Lemma 1, $\mathbb{O} = \sum_{\mathcal{S}_1} \alpha_i + \sum_{\mathcal{S}_1} \sigma_i^U + F(\mathcal{S}_2) + \theta_{\mathcal{S}_2}^U$.

Lemmas 2 and 4 are used to determine the range around kT in which each of the stable queues fluctuates.

Lemma 2: Suppose that there are numbers $H \geq 0$, $\Phi > 0$, and queue sets $\mathcal{L} \subset \mathcal{N}$, $\mathcal{D} \subseteq \mathcal{N} - \mathcal{L}$, such that the following hold.

a) The set $\mathcal{G}(t) = \{i \in \mathcal{D} : q_i(t) > H\}$, has highest priority at time t among the queues in $\mathcal{N} - \mathcal{L}$.

b) If $\max_{i \in \mathcal{D}}\{q_i(t)\} \leq H + \Phi$, the queues in the set $\mathcal{G}(t)$ are served according to the LCQ policy.

c) For any $\mathcal{S} \subseteq \mathcal{D}$ it holds

$$\sum_{\mathcal{S}} \alpha_i \leq F(\mathcal{L} \cup \mathcal{S}) - F(\mathcal{L}). \tag{12}$$

d) $q_i(0) \leq H$ for all $i \in \mathcal{D}$,

then, there is a number $\mathbb{O}$ such that, if $\Phi > \mathbb{O}$, it holds

$$\max_{i \in \mathcal{D}}\{q_i(t)\} \leq H + \mathbb{O}, \text{ for all } t \geq 0.$$

Proof: Let $D = |\mathcal{D}|$. Also define

$$y_l(t) = \sum_{\mathcal{S}_{\mathcal{D}}^{(l)}(t)} (q_i(t) - H)^+.$$

We observe that for $2 \leq l \leq D - 1$, it holds

$$y_{l-1}(t) + y_{l+1}(t) =$$
$$2y_l(t) + \left(q_{\mathcal{D}}^{(l+1)}(t) - H\right)^+ - \left(q_{\mathcal{D}}^{(l)}(t) - H\right)^+, \tag{13}$$

and

$$y_2(t) = 2y_1(t)$$
$$+ \left(q_{\mathcal{D}}^{(2)}(t) - H\right)^+ - \left(q_{\mathcal{D}}^{(1)}(t) - H\right)^+, \tag{14}$$

$$y_{D-1}(t) = y_D(t) - \left(q_{\mathcal{D}}^{(D)}(t) - H\right)^+. \tag{15}$$

Consider a time t where $q_{\mathcal{D}}^{(l)}(t) > H$ and let $t_0 - 1$ be the largest time before t, such that $\mathcal{S}_{\mathcal{D}}^{(l)}(t) \neq \mathcal{S}_{\mathcal{D}}^{(l)}(t_0 - 1)$ or $q_{\mathcal{D}}^{(l)}(t_0 - 1) \leq H$ (for $l = D$ only the second situation makes sense). Time t_0 is well defined because of assumption d) of the lemma. In the time interval $[t_0, t]$, the set $\mathcal{S}_{\mathcal{D}}^{(l)}(t)$ remains the same and all the queues in this set are bigger than H, i.e., nonempty. Provided that Φ is large enough (as will be seen it suffices to be $\mathbb{O}$), this set of queues has priority over the queues in $\mathcal{N} - \mathcal{L}$ and uses all the available slots in $[t_0, t]$. Since these slots are at least $C_{\mathcal{L} \cup \mathcal{S}_{\mathcal{D}}^{(l)}(t_0)}(t_0, t) - C_{\mathcal{L}}(t_0, t)$, we have $B_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)}(t_0, t) \geq C_{\mathcal{L} \cup \mathcal{S}_{\mathcal{D}}^{(l)}(t_0)}(t_0, t) - C_{\mathcal{L}}(t_0, t)$. Setting $\overline{F}(\mathcal{S}) = F(\mathcal{L} \cup \mathcal{S}) - F(\mathcal{L})$, we conclude

$$\sum_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)} q_i(t) = \sum_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)} q_i(t_0) + A_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)}(t_0, t)$$
$$- B_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)}(t_0, t)$$
$$\leq \sum_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)} q_i(t_0)$$
$$+ \left(\sum_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)} \alpha_i - \overline{F}\left(\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)\right)\right)(t - t_0) + \mathbb{O}$$
$$\leq \sum_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)} q_i(t_0) + \mathbb{O}. \tag{16}$$

In the first inequality above we used the burstiness constraints on the arrival and channel availability processes. In the second inequality we used assumption c) of the lemma. Because of the way t_0 is defined, we have that for $i \in \mathcal{S}_{\mathcal{D}}^{(l)}(t_0)$ it holds $q_i(t) > H$ and $q_i(t_0) > H$. Therefore by subtracting H, l times from both sides of inequality (16) we obtain

$$\sum_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)} (q_i(t) - H)^+$$
$$\leq \sum_{\mathcal{S}_{\mathcal{D}}^{(l)}(t_0)} (q_i(t_0) - H)^+ + \mathbb{O}. \tag{17}$$

Hence,

$$y_l(t) \leq y_l(t_0) + \mathbb{O}. \tag{18}$$

Let $2 \leq l \leq D - 1$. We will show now that

$$y_l(t_0) \leq \frac{1}{2}y_{l-1}(t_0) + \frac{1}{2}y_{l+1}(t_0) + \mathbb{O}. \tag{19}$$

Note that by definition $q_{\mathcal{D}}^{(l)}(t_0) > H$ and therefore from equation (13) it holds

$$y_l(t_0) = \frac{1}{2}y_{l-1}(t_0) + \frac{1}{2}y_{l+1}(t_0) + \frac{1}{2}\left(q_{\mathcal{D}}^{(l)}(t_0) - H\right)$$
$$- \frac{1}{2}\left(q_{\mathcal{D}}^{(l+1)}(t_0) - H\right)^+. \tag{20}$$

We consider two cases.

Case 1. $q_{\mathcal{D}}^{(l)}(t_0 - 1) \leq H$. Then, there must be an index $i_0 \in \mathcal{S}_{\mathcal{D}}^{(l)}(t_0)$ such that $q_{i_0}(t_0 - 1) \leq H$ and $q_{\mathcal{D}}^{(l)}(t_0) \leq q_{i_0}(t_0)$. Using Lemma 1 with $\mathcal{S}_1 = \{i_0\}$, $\mathcal{S}_2 = \varnothing$, we conclude that

$$q_{i_0}(t_0) \leq q_{i_0}(t_0 - 1) + \mathbb{O}$$
$$\leq H + \mathbb{O}.$$

Hence,

$$q_{\mathcal{D}}^{(l)}(t_0) - H \leq \mathbb{O}. \qquad (21)$$

From (21) and (20) (taking also into account that $\left(q_{\mathcal{D}}^{(l+1)}(t_0) - H\right)^+ \geq 0$) follows (19).

Case 2. $\mathcal{S}_{\mathcal{D}}^{(l)}(t_0) \neq \mathcal{S}_{\mathcal{D}}^{(l)}(t_0 - 1)$ and $q_{\mathcal{D}}^{(l)}(t_0 - 1) > H$. Then, there must be indices $i_0 \in \mathcal{S}_{\mathcal{D}}^{(l)}(t_0)$, $j_0 \in \overline{\mathcal{S}}_{\mathcal{D}}^{(l)}(t_0)$, such that $q_{i_0}(t_0 - 1) \leq q_{j_0}(t_0 - 1)$ and $q_{\mathcal{D}}^{(l)}(t_0) - q_{\mathcal{D}}^{(l+1)}(t_0) \leq q_{i_0}(t_0) - q_{j_0}(t_0)$. By using Lemma 1 with $\mathcal{S}_1 = \{i_0\}$, $\mathcal{S}_2 = \{j_0\}$, we conclude that

$$q_{i_0}(t_0) - q_{j_0}(t_0) \leq q_{i_0}(t_0 - 1) - q_{j_0}(t_0 - 1) + \mathbb{O} \leq \mathbb{O}.$$

Hence,

$$q_{\mathcal{D}}^{(l)}(t_0) - q_{\mathcal{D}}^{(l+1)}(t_0) \leq \mathbb{O}. \qquad (22)$$

Since $\left(q_{\mathcal{D}}^{(l+1)}(t_0) - H\right)^+ \geq q_{\mathcal{D}}^{(l+1)}(t_0) - H$, from (20) we have

$$y_l(t_0) \leq \frac{1}{2} y_{l-1}(t_0) + \frac{1}{2} y_{l+1}(t_0)$$
$$+ \frac{1}{2}\left(q_{\mathcal{D}}^{(l)}(t_0) - q_{\mathcal{D}}^{(l+1)}(t_0)\right).$$

which in conjunction with (22) shows (19).

Similarly, we have with an analogous definition of t_0,

$$y_1(t) \leq \frac{1}{2} y_2(t_0) + \mathbb{O}, \qquad (23)$$

$$y_D(t) \leq y_{D-1}(t_0) + \mathbb{O}. \qquad (24)$$

If $q_{\mathcal{D}}^{(l)}(t) \leq H$, then we have from (13), (14) and (15) that (19) as well as (23) and (24) still hold with $t_0 = t$.

Fix now a time $\overline{t}$ and define

$$\overline{y}_l(\overline{t}) = \max_{t \leq \overline{t}} y_l(t) < \infty. \qquad (25)$$

From (18), (19), (24) and (23) it holds for $2 \leq l \leq D - 1$, that for any t, $0 \leq t \leq \overline{t}$,

$$y_l(t) \leq \frac{1}{2} \overline{y}_{l-1}(\overline{t}) + \frac{1}{2} \overline{y}_{l+1}(\overline{t}) + \mathbb{O}, \qquad (26)$$

and

$$y_1(t) \leq \frac{1}{2} \overline{y}_2(\overline{t}) + \mathbb{O} \qquad (27)$$

$$y_D(t) \leq \overline{y}_{D-1}(\overline{t}) + \mathbb{O} \qquad (28)$$

Therefore we have for $2 \leq l \leq D - 1$,

$$\overline{y}_l(\overline{t}) \leq \frac{1}{2} \overline{y}_{l-1}(\overline{t}) + \frac{1}{2} \overline{y}_{l+1}(\overline{t}) + \mathbb{O}, \qquad (29)$$

and

$$\overline{y}_1(\overline{t}) \leq \frac{1}{2} \overline{y}_2(\overline{t}) + \mathbb{O}, \qquad (30)$$

$$\overline{y}_D(\overline{t}) \leq \overline{y}_{D-1}(\overline{t}) + \mathbb{O}. \qquad (31)$$

The above inequalities can be written in matrix form as:

$$(\mathbf{I} - \mathbf{B})\mathbf{Y} \leq \overline{\mathbb{O}}, \qquad (32)$$

where $\mathbf{Y} = \left[\overline{y}_1(\overline{t})...\overline{y}_D(\overline{t})\right]^T$, and $\mathbf{I}$ is the unity matrix, $\overline{\mathbb{O}}$ is a matrix whose elements are of type $\mathbb{O}$ and

$$\mathbf{B} = \begin{bmatrix} 0 & 1/2 & 0 & 0 & 0 \\ 1/2 & 0 & 1/2 & 0 & 0 \\ 0 & 1/2 & 0 & 1/2 & 0 \\ ... & ... & ... & ... & ... \\ 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}.$$

Since the row sums of $\mathbf{B}$ are all less than or equal to 1 and the sum of the first raw is $1/2$, i.e., less than 1, it follows from the Perron-Frobenius Theorem [10], that the eigenvalues of $\mathbf{B}$ are all smaller than 1 in absolute value. Therefore, the matrix $(\mathbf{I} - \mathbf{B})^{-1}$ has nonnegative elements. Hence, we can multiply (32) with $(\mathbf{I} - \mathbf{B})^{-1}$ to get $\mathbf{Y} \leq (\mathbf{I} - \mathbf{B})^{-1}\mathbb{O}$, or

$$\overline{y}_l(\overline{t}) \leq \mathbb{O}, \ l = 1, ..., D. \qquad (33)$$

Therefore,

$$\lim_{\overline{t} \to \infty} \overline{y}_l(\overline{t}) = \sup_t y_l(t) = y_l \leq \mathbb{O}, \ l = 1, ..., D. \qquad (34)$$

Hence

$$\max_{i \in \mathcal{D}} \{q_i(t)\} - H = q_{\mathcal{D}}^{(1)}(t) - H$$
$$\leq \left(q_{\mathcal{D}}^{(1)}(t) - H\right)^+ = y_1(t) \leq y_1 \leq \mathbb{O}$$

and the lemma follows. ∎

The next lemma shows that when the inequalities in (12) are strengthened to strict inequalities, then the queues sizes are bounded by $H + \mathbb{O}$ after some time slot, under any finite initial condition on the queue sizes.

Lemma 3: Suppose that conditions a) and b) of Lemma 2 hold and in addition, for any $\mathcal{S} \subseteq \mathcal{D}$, it holds

$$\sum_{\mathcal{S}} \alpha_i < F(\mathcal{L} \cup \mathcal{S}) - F(\mathcal{L}) .$$

Then under any initial queue sizes $q_i(0) < \infty$, $i \in \mathcal{D}$, there is a finite time τ_0 such that

$$\max_{i \in \mathcal{D}} \{q_i(t)\} \leq H + \mathbb{O}, \ \text{for all } t \geq \tau_0.$$

Proof: The proof is omitted. ∎

The next lemma provides conditions under which it is known that the queue sizes of certain set of queues do not fall below certain threshold after some time slot. Its proof is analogous to the proof of Lemma 3.

Lemma 4: If there is a positive number H and queue sets $\mathcal{L} \subset \mathcal{N}$, $\mathcal{D} \subseteq \mathcal{N} - \mathcal{L}$, such that the following hold.

a) The queues in $\mathcal{L}$ always have packets to transmit and have higher priority than the queues in $\mathcal{G}(t) = \{i \in \mathcal{D} : q_i(t) \leq H\}$.

b) The queues in the set $\mathcal{G}(t)$, are served according to LCQ policy.

c) For any subset $\mathcal{S} \subseteq \mathcal{D}$, with $\overline{\mathcal{S}} = \mathcal{D} - \mathcal{S}$ it holds

$$\sum_{\mathcal{S}} \alpha_i > F(\mathcal{L} \cup \mathcal{D}) - F(\mathcal{L} \cup \overline{\mathcal{S}}),$$

then there it a time τ_0 such that $\min_{i \in \mathcal{D}} \{q_i(t)\} \geq H - \mathbb{O}$, for all $t \geq \tau_0$.

Fig. 2. Partitioning of the index set.

Next we need to examine in more detail the structure of the optimal linear programming solution (11). According to (11), x_k^* may take values less than or equal to α_k. An index such that $x_k^* = \alpha_k$ is called "stable", while an index such that $x_k^* < \alpha_k$, "unstable". We therefore have that for a stable index k, for any $\mathcal{D} \subseteq \mathcal{N}_{k-1}$

$$\alpha_k \leq F(k \cup \mathcal{D}) - \sum_{\mathcal{D}} x_i^*. \tag{35}$$

Similarly, for an unstable index k, for any set $\mathcal{D}_k \subseteq \mathcal{N}_{k-1}$ such that

$$x_k^* = F(k \cup \mathcal{D}_k) - \sum_{\mathcal{D}_k} x_i^*,$$

it holds

$$F(k \cup \mathcal{D}_k) - \sum_{\mathcal{D}_k} x_i^* < \alpha_k. \tag{36}$$

The general structure of the vector $\{x_k^*\}_{k \in \mathcal{N}}$ is as follows. The set of indices is partitioned into index sets $\mathcal{I}_i^s$, $i = 1, ..., l_s$, $\mathcal{I}_i^u$, $i = 1, ..., l_u$ such that

- Indices in the set $\cup_{i=1}^{l_s} \mathcal{I}_i^s$ are stable. Indices in the set $\cup_{i=1}^{l_u} \mathcal{I}_i^u$ are unstable.
- Index set $\mathcal{I}_i^x$, $x \in \{s, u\}$ consists of successive integers.
- If $i > j$ then all indices in $\mathcal{I}_i^x$, $x \in \{s, u\}$ are larger than the indices in $\mathcal{I}_j^x$. Figure 2 shows an example of the partition of the index set for $N = 10$ channels. For convenience in the discussion we assume that for a given i, the indices in $\mathcal{I}_i^s$ are smaller than the indices in $\mathcal{I}_i^u$. Hence, for consistency, if index 1 is unstable, we define $\mathcal{I}_1^s = \varnothing$.

If $k \in \mathcal{I}_j^u$ then define $\widehat{\mathcal{I}}_k = \cup_{i=1}^j \mathcal{I}_i^s$.

Denote by $u_1 < u_2 < ... < u_L$, $L = \left| \cup_{i=1}^{l_u} \mathcal{I}_i^u \right|$ the unstable indices.

In the following, we assume that strict submodularity holds, i.e., whenever $\mathcal{S} - \mathcal{D} \neq \varnothing$ and $\mathcal{D} - \mathcal{S} \neq \varnothing$, it holds,

$$F(\mathcal{S}) + F(\mathcal{D}) > F(\mathcal{S} \cup \mathcal{D}) + F(\mathcal{S} \cap \mathcal{D}), \tag{37}$$

therefore, whenever $F(\mathcal{S}) + F(\mathcal{D}) = F(\mathcal{S} \cup \mathcal{D}) + F(\mathcal{S} \cap \mathcal{D})$, then either $\mathcal{S} \subseteq \mathcal{D}$ or $\mathcal{D} \subseteq \mathcal{S}$. This assumption is not essential but simplifies the proofs.

Lemma 5: Consider the vector $\{x_k^*\}_{k \in \mathcal{N}}$ defined for $k = 1, ..., N$, by the recursion

$$x_k^* = \min \left\{ \alpha_k, \min_{\mathcal{D} \subseteq \mathcal{N}_{k-1}} \left\{ F(k \cup \mathcal{D}) - \sum_{\mathcal{D}} x_i^* \right\} \right\}.$$

a) For any set $\mathcal{S} \subseteq \mathcal{I}_1^s$ it holds

$$\sum_{\mathcal{S}} \alpha_i \leq F(\mathcal{S}).$$

b) With index $k \in \mathcal{I}_j^u$, $j = 1, ..., l_u$, there is an associated index set $\widehat{\mathcal{D}}_k$ such that $\widehat{\mathcal{D}}_k \subseteq \mathcal{N}_{k-1}$,

$$\sum_{k \cup \widehat{\mathcal{D}}_k} x_i^* = F\left(k \cup \widehat{\mathcal{D}}_k \right),$$

for all $\mathcal{D} \subseteq \mathcal{N}_{k-1}$

$$F\left(k \cup \widehat{\mathcal{D}}_k \right) - \sum_{\widehat{\mathcal{D}}_k} x_i^* \leq F(k \cup \mathcal{D}) - \sum_{\mathcal{D}} x_i^*,$$

and for all $\mathcal{D} \subset \widehat{\mathcal{D}}_k$,

$$x_k^* = F\left(k \cup \widehat{\mathcal{D}}_k \right) - \sum_{\widehat{\mathcal{D}}_k} x_i^* < F(k \cup \mathcal{D}) - \sum_{\mathcal{D}} x_i^*.$$

c) $u_{i-1} \cup \widehat{\mathcal{D}}_{u_{i-1}} \subseteq \widehat{\mathcal{D}}_{u_i}$, $2 \leq i \leq L$.
d) $\widehat{\mathcal{D}}_{u_i} = \cup_{j=1}^{i-1} u_j \cup \widetilde{\mathcal{D}}_{u_i}$, where $\widetilde{\mathcal{D}}_{u_i} \subseteq \widehat{\mathcal{I}}_{u_i}$, $1 \leq i \leq L$.
e) If k is an unstable index, then for any set $\mathcal{S} \subseteq \mathcal{H}_k \equiv \widehat{\mathcal{I}}_k - \widetilde{\mathcal{D}}_k$ it holds

$$\sum_{\mathcal{S}} \alpha_i \leq F\left(k \cup \widehat{\mathcal{D}}_k \cup \mathcal{S} \right) - F\left(k \cup \widehat{\mathcal{D}}_k \right).$$

f) If k is the largest index in $\mathcal{I}_j^u$ then for any set $\mathcal{S} \subseteq \left(\widehat{\mathcal{I}}_k \cup \mathcal{I}_{j+1}^s \right) - \widetilde{\mathcal{D}}_k$ it holds

$$\sum_{\mathcal{S}} \alpha_i \leq F\left(k \cup \widehat{\mathcal{D}}_k \cup \mathcal{S} \right) - F\left(k \cup \widehat{\mathcal{D}}_k \right).$$

Proof: a)Let $\mathcal{S} \subseteq \mathcal{I}_1^s$ and let k be the largest index in $\mathcal{S}$. Then, since $\mathcal{S} - \{k\} \in \mathcal{N}_{k-1}$, according to the definition of x_k^* and $\mathcal{I}_1^s$ we have

$$\alpha_k \leq F(k \cup (\mathcal{S} - k)) - \sum_{\mathcal{S} - k} \alpha_i$$

or,

$$\sum_{\mathcal{S}} \alpha_i \leq F(\mathcal{S}). \tag{38}$$

b) According to the definition of x_k^*, there is a set $\mathcal{D}_k \subseteq \mathcal{N}_{k-1}$ such that for all $\mathcal{D} \subseteq \mathcal{N}_{k-1}$,

$$x_k^* = F(k \cup \mathcal{D}_k) - \sum_{\mathcal{D}_k} x_i^*$$

$$\leq F(k \cup \mathcal{D}) - \sum_{\mathcal{D}} x_i^*. \tag{39}$$

Let $\mathcal{D}_k^1 \subseteq \mathcal{N}_{k-1}$ be another set such that

$$x_k^* = F\left(k \cup \mathcal{D}_k^1 \right) - \sum_{\mathcal{D}_k^1} x_i^*.$$

Since by definition

$$F\left(k \cup \mathcal{D}_k^1\right) - \sum_{\mathcal{D}_k^1} x_i^* \leq F(k \cup \mathcal{D}_k^1 \cup \mathcal{D}_k) - \sum_{\mathcal{D}_k^1 \cup \mathcal{D}_k} x_i^*,$$

$$F\left(k \cup \mathcal{D}_k\right) - \sum_{\mathcal{D}_k} x_i^* \leq F(k \cup (\mathcal{D}_k^1 \cap \mathcal{D}_k)) - \sum_{\mathcal{D}_k^1 \cap \mathcal{D}_k} x_i^*,$$

we have by adding and cancelling summation terms,

$$F\left(k \cup \mathcal{D}_k\right) + F(k \cup \mathcal{D}_k^1) \leq$$
$$F(k \cup \mathcal{D}_k^1 \cup \mathcal{D}_k) + F(k \cup (\mathcal{D}_k^1 \cap \mathcal{D}_k)) =$$
$$F(k \cup \mathcal{D}_k^1 \cup \mathcal{D}_k) + F((k \cup \mathcal{D}_k^1) \cap (k \cup \mathcal{D}_k)).$$

By submodularity,

$$F\left(k \cup \mathcal{D}_k\right) + F\left(k \cup \mathcal{D}_k^1\right)$$
$$= F(k \cup \mathcal{D}_k^1 \cup \mathcal{D}_k) + F((k \cup \mathcal{D}_k^1) \cap (k \cup \mathcal{D}_k)).$$

This equality and the strict submodularity property imply that either $\mathcal{D}_k \subseteq \mathcal{D}_k^1$ or $\mathcal{D}_k^1 \subseteq \mathcal{D}_k$. Let now $\widehat{\mathcal{D}}_k$ be the set with smallest cardinality among those that satisfy (39). By definition, this set satisfies part b) of the lemma.

Assume for the moment that part c) holds up to value i. Part d) for $i = 1$ follows from the definitions, and for $i \geq 2$ it is immediate from c). We will also show that e) and f) hold.

To prove e), let $\mathcal{S} \subseteq \widehat{\mathcal{I}}_{u_i} - \widehat{\mathcal{D}}_{u_i} = \widehat{\mathcal{I}}_{u_i} - \widehat{\mathcal{D}}_{u_i} \subseteq \mathcal{N}_{u_i - 1}$. Then,

$$F(u_i \cup \widehat{\mathcal{D}}_{u_i}) - \sum_{\widehat{\mathcal{D}}_{u_i}} x_j^* \leq F(u_i \cup \widehat{\mathcal{D}}_{u_i} \cup \mathcal{S}) - \sum_{\widehat{\mathcal{D}}_{u_i} \cup \mathcal{S}} x_j^*,$$

hence, taking into account that $\mathcal{S} \cap \widehat{\mathcal{D}}_{u_i} = \varnothing$,

$$\sum_{\mathcal{S}} x_j^* \leq F(u_i \cup \widehat{\mathcal{D}}_{u_i} \cup \mathcal{S}) - F(u_i \cup \widehat{\mathcal{D}}_{u_i}). \quad (40)$$

To prove f) assume that u_i is the largest index in $\mathcal{I}_j^u$ and $\mathcal{I}_{j+1}^s \neq \varnothing$. Let $\mathcal{S} \subseteq \left(\widehat{\mathcal{I}}_{u_i} \cup \mathcal{I}_{j+1}^s\right) - (u_i \cup \widehat{\mathcal{D}}_{u_i})$ and let i_k be the largest index in $\mathcal{S}$. If $i_k \in \mathcal{I}_j^s$ then $\mathcal{S} \subseteq \widehat{\mathcal{I}}_{u_i} - \widehat{\mathcal{D}}_{u_i}$ and (40) holds. If $i_k \in \mathcal{I}_{j+1}^s$, then since $u_i \cup \widehat{\mathcal{D}}_{u_i} \cup (\mathcal{S} - i_k) \subseteq \mathcal{N}_{i_k - 1}$, and $\left(u_i \cup \widehat{\mathcal{D}}_{u_i}\right) \cap (\mathcal{S} - i_k) = \varnothing$, we have from (35),

$$\alpha_{i_k} \leq F(i_k \cup u_i \cup \widehat{\mathcal{D}}_{u_i} \cup (\mathcal{S} - i_k)) - \sum_{u_i \cup \widehat{\mathcal{D}}_{u_i}} x_j^* - \sum_{\mathcal{S} - i_k} \alpha_j$$
$$= F(u_i \cup \widehat{\mathcal{D}}_{u_i} \cup \mathcal{S}) - F(u_i \cup \widehat{\mathcal{D}}_{u_i}) - \sum_{\mathcal{S} - i_k} \alpha_j,$$

where we used the fact that by part b) of the lemma,

$$\sum_{u_i \cup \widehat{\mathcal{D}}_{u_i}} x_j^* = F(u_i \cup \widehat{\mathcal{D}}_{u_i}).$$

That is,

$$\sum_{\mathcal{S}} \alpha_j \leq F(u_i \cup \widehat{\mathcal{D}}_{u_i} \cup \mathcal{S}) - F(u_i \cup \widehat{\mathcal{D}}_{u_i}).$$

It remains to prove part c). Assume that part c) holds up to index $i - 1$. Assume also that for index u_{m-1}, $1 < m < i$ it holds $u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}} \subseteq \widehat{\mathcal{D}}_{u_i}$. We will show then that it also

holds, $u_m \cup \widehat{\mathcal{D}}_{u_m} \subseteq \widehat{\mathcal{D}}_{u_i}$, which will prove part c), provided that we also show that

$$u_1 \cup \widehat{\mathcal{D}}_{u_1} \subseteq \widehat{\mathcal{D}}_{u_i}. \quad (41)$$

The proof of (41) will be outlined at the end, since it is basically a rewording of the argument for general m.

Since $\widehat{\mathcal{D}}_{u_i} \cup u_m \cup \widehat{\mathcal{D}}_{u_m} \subseteq \mathcal{N}_{u_i - 1}$,

$$F\left(u_i \cup \widehat{\mathcal{D}}_{u_i}\right) - \sum_{\widehat{\mathcal{D}}_{u_i}} x_j^*$$

$$\leq F\left(u_i \cup \widehat{\mathcal{D}}_{u_i} \cup u_m \cup \widehat{\mathcal{D}}_{u_m}\right) - \sum_{\widehat{\mathcal{D}}_{u_i} \cup u_m \cup \widehat{\mathcal{D}}_{u_m}} x_j^*$$

$$= F\left(u_i \cup \widehat{\mathcal{D}}_{u_i} \cup u_m \cup \widehat{\mathcal{D}}_{u_m}\right) - \sum_{u_m \cup \widehat{\mathcal{D}}_{u_m}} x_j^*$$

$$- \sum_{\widehat{\mathcal{D}}_{u_i} - (u_m \cup \widehat{\mathcal{D}}_{u_m})} x_j^*$$

$$= F\left(u_i \cup \widehat{\mathcal{D}}_{u_i} \cup u_m \cup \widehat{\mathcal{D}}_{u_m}\right) - F(u_m \cup \widehat{\mathcal{D}}_{u_m})$$

$$- \sum_{\widehat{\mathcal{D}}_{u_i} - (u_m \cup \widehat{\mathcal{D}}_{u_m})} x_j^*.$$

That is

$$\sum_{\widehat{\mathcal{D}}_{u_i} \cap (u_m \cup \widehat{\mathcal{D}}_{u_m})} x_j^*$$
$$\geq F(u_i \cup \widehat{\mathcal{D}}_{u_i}) + F(u_m \cup \widehat{\mathcal{D}}_{u_m})$$
$$- F(u_i \cup \widehat{\mathcal{D}}_{u_i} \cup u_m \cup \widehat{\mathcal{D}}_{u_m})$$
$$\geq F\left(\left(u_i \cup \widehat{\mathcal{D}}_{u_i}\right) \cap \left(u_m \cup \widehat{\mathcal{D}}_{u_m}\right)\right) \quad (42)$$
$$= F\left(\widehat{\mathcal{D}}_{u_i} \cap \left(u_m \cup \widehat{\mathcal{D}}_{u_m}\right)\right), \quad (43)$$

where in (42) the inequality is strict if $\left(u_i \cup \widehat{\mathcal{D}}_{u_i}\right) - \left(u_m \cup \widehat{\mathcal{D}}_{u_m}\right) \neq \varnothing$ and $\left(u_m \cup \widehat{\mathcal{D}}_{u_m}\right) - \left(u_i \cup \widehat{\mathcal{D}}_{u_i}\right) \neq \varnothing$. Equality (43) follows from the fact that since $i > m$, $u_i \cap \left(u_m \cup \widehat{\mathcal{D}}_{u_m}\right) = \varnothing$.

We also have since $\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m} \subseteq \mathcal{N}_{u_m - 1}$,

$$F\left(u_m \cup \widehat{\mathcal{D}}_{u_m}\right) - \sum_{\widehat{\mathcal{D}}_{u_m}} x_j^*$$

$$\leq F\left(u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)\right) - \sum_{\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}} x_j^*$$

$$= F\left(u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)\right) - \sum_{u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)} x_j^* + x_{u_m}^*,$$

where in the last equality we used the fact that $u_m \notin \widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}$. Since by part b), $x_{u_m}^* = F\left(u_m \cup \widehat{\mathcal{D}}_{u_m}\right) - \sum_{\widehat{\mathcal{D}}_{u_m}} x_j^*$, we conclude that

$$\sum_{u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)} x_j^* \leq F\left(u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)\right), \quad (44)$$

where in (44), by part b) of the lemma, the inequality is strict if $\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m} \subset \widehat{\mathcal{D}}_{u_m}$.

Assume now that $u_m \notin \widehat{\mathcal{D}}_{u_i}$. Then $\widehat{\mathcal{D}}_{u_i} \cap \left(u_m \cup \widehat{\mathcal{D}}_{u_m}\right) = \widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}$. Define $\widehat{\mathcal{E}}_m = \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right) - \left(u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}}\right)$. Since by the inductive assumption we have $u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}} \subseteq \widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}$, we can write

$$\sum_{\widehat{\mathcal{D}}_{u_i} \cap (u_m \cup \widehat{\mathcal{D}}_{u_m})} x_j^* = \sum_{\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}} x_j^*$$
$$= \sum_{\widehat{\mathcal{E}}_m} x_j^* + \sum_{u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}}} x_j^*$$
$$= \sum_{\widehat{\mathcal{E}}_m} x_j^* + F\left(u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}}\right).$$

Since $u_m \notin \widehat{\mathcal{D}}_{u_i}$ and $u_i \notin \widehat{\mathcal{D}}_{u_m}$, from the strict submodularity and (43) it follows that

$$\sum_{\widehat{\mathcal{E}}_m} x_j^* > F\left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)$$
$$- F\left(u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}}\right). \quad (45)$$

Observe that by the inductive hypothesis we have that $\widehat{\mathcal{E}}_m \subseteq \widehat{\mathcal{I}}_{u_{m-1}} - \widetilde{\mathcal{D}}_{u_{m-1}}$ or $\widehat{\mathcal{E}}_m \subseteq \left(\widehat{\mathcal{I}}_{u_{m-1}} \cup \mathcal{I}_{j+1}^s\right) - \widetilde{\mathcal{D}}_{u_{m-1}}$ depending on whether u_{m-1} is not, or is, the largest element in $\mathcal{I}_j^u$. In either case have by e) and f)

$$\sum_{\widehat{\mathcal{E}}_m} x_j^* = \sum_{\widehat{\mathcal{E}}_m} \alpha_j \leq F\left(u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}} \cup \widehat{\mathcal{E}}_m\right)$$
$$- F\left(u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}}\right)$$
$$= F\left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right) - F\left(u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}}\right). \quad (46)$$

Inequality (45) contradicts (46) and therefore we conclude that $u_m \in \widehat{\mathcal{D}}_{u_i}$. This implies that

$$\widehat{\mathcal{D}}_{u_i} \cap (u_m \cup \widehat{\mathcal{D}}_{u_m}) = u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)$$

From (43) and (44) it follows that

$$F\left(u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)\right) \geq \sum_{u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)} x_j^*$$
$$\geq F\left(u_m \cup \left(\widehat{\mathcal{D}}_{u_i} \cap \widehat{\mathcal{D}}_{u_m}\right)\right),$$

with equalities holding only if $\widehat{\mathcal{D}}_{u_m} \subseteq \widehat{\mathcal{D}}_{u_i}$.

To complete the proof we need to prove (41). For this we follow essentially the same arguments as above, by replacing m with 1 and $u_{m-1} \cup \widehat{\mathcal{D}}_{u_{m-1}}$ with the empty set. ∎

According to parts b) and c) of Lemma 5, $\widehat{\mathcal{D}}_{u_i}$ can be written as $\widehat{\mathcal{D}}_{u_i} = \cup_{j=1}^{i-1} u_j \cup \widetilde{\mathcal{D}}_{u_i}$, where $\widetilde{\mathcal{D}}_{u_i} \in \mathcal{I}_{u_i}$. Let $\mathcal{D} \subset \widetilde{\mathcal{D}}_{u_i}$, hence $(\cup_{j=1}^{i-1} u_j \cup \mathcal{D}) \subset \widehat{\mathcal{D}}_{u_i}$. Applying part b) of Lemma 5 and taking into account that $x_i^* = \alpha_i$ when $i \in \widetilde{\mathcal{D}}_{u_i}$, we have

$$F\left(\cup_{j=1}^i u_j \cup \widetilde{\mathcal{D}}_{u_i}\right) - \sum_{\cup_{l=1}^{i-1} u_l} x_j^* - \sum_{\widetilde{\mathcal{D}}_{u_i}} \alpha_j$$
$$< F\left(\cup_{j=1}^i u_j \cup \mathcal{D}\right) - \sum_{\cup_{l=1}^{i-1} u_l} x_j^* - \sum_{\mathcal{D}} \alpha_j,$$

or

$$F\left(\cup_{j=1}^i u_j \cup \widetilde{\mathcal{D}}_{u_i}\right) - F\left(\cup_{j=1}^i u_j \cup \mathcal{D}\right) < \sum_{\widetilde{\mathcal{D}}_{u_i} - \mathcal{D}} \alpha_i.$$

Setting $\mathcal{S} = \widetilde{\mathcal{D}}_{u_i} - \mathcal{D}$ in the previous inequality we have the following corollary.

Corollary 6: If we set $\mathcal{U}_i = \cup_{j=1}^i u_j$, then for any $\mathcal{S} \subseteq \widetilde{\mathcal{D}}_{u_i}$, $\mathcal{S} \neq \varnothing$, it holds

$$F\left(\mathcal{U}_i \cup \widetilde{\mathcal{D}}_{u_i}\right) - F\left(\mathcal{U}_i \cup \left(\widetilde{\mathcal{D}}_{u_i} - \mathcal{S}\right)\right) < \sum_{\mathcal{S}} \alpha_i.$$

The next theorem provides the range within which each of the queues in $\mathcal{N}$ fluctuates.

Theorem 7: Under policy π^*, if $q_i(0) = 0$, $i \in \mathcal{N}$, then for T large enough, the queues in $\cup_{i=1}^L \mathcal{I}_i^u$ tend to infinity. Moreover, setting $\widetilde{\mathcal{D}}_{u_0} = \varnothing$, we have for a number $M = \mathbb{O}$,

$$\max_{\widetilde{\mathcal{D}}_{u_j} - \widetilde{\mathcal{D}}_{u_{j-1}}} \{q_i(t)\} \leq (N + 1 - u_j)T + M \text{ for all } t \geq 0,$$

$$\min_{\widetilde{\mathcal{D}}_{u_j} - \widetilde{\mathcal{D}}_{u_{j-1}}} \{q_i(t)\} \geq (N + 1 - u_j)T - M \text{ for all } t \geq \tau_0,$$

and if $i \notin \widetilde{\mathcal{D}}_{u_L}$, $\max_{\cup_{j=1}^{l_s} \mathcal{I}_j^s - \widetilde{\mathcal{D}}_{u_L}} q_i(t) \leq M$ for all t. It suffices to take $T > 2M$.

In the course of the proof of Theorem 7 we also prove the main result of this paper, that is,

Theorem 8: The throughputs achieved by policy π^* for T large enough, satisfy the resurcive equations

$$x_k^* = \min\left\{\alpha_k, \min_{\mathcal{D} \subseteq \mathcal{N}_{k-1}} \left\{F\left(k \cup \mathcal{D}\right) - \sum_{\mathcal{D}} x_i^*\right\}\right\},$$

for $k = 1, ..., N$.

Proof: In Lemma 2 set $\mathcal{D} = \mathcal{I}_1^s$, $H = (N + 1 - u_1)T$, $\Phi = T$, and $\mathcal{L} = \varnothing$. Conditions a), b) of the lemma hold because of the definition of policy π^*. From Lemma 5 we have that condition c) holds as well. Since condition d) of the lemma also holds by the assumption that $q_i(0) = 0$, $i \in \mathcal{N}$, we conclude that for all $t \geq 0$,

$$\max_{i \in \mathcal{I}_1^s} \{q_i(t)\} \leq (N + 1 - u_1)T + \mathbb{O} \quad (47)$$

provided that $T \geq \mathbb{O}$.

Since $B_i(0, t) = A_i(0, t) - q_i(t)$, from (47) we conclude that for $i \in \mathcal{I}_1^s$,

$$r_i = \lim_{t \to \infty} \frac{B_i(0, t)}{t} = \lim_{t \to \infty} \frac{A_i(0, t)}{t} = \alpha_i = x_i^*. \quad (48)$$

Since $\widehat{\mathcal{D}}_{u_1} = \widetilde{\mathcal{D}}_{u_1} \subseteq \mathcal{I}_1^s$ we have from (36) and the definition of $\mathcal{I}_1^s$ that

$$F(u_1 \cup \widetilde{\mathcal{D}}_{u_1}) - \sum_{\widetilde{\mathcal{D}}_{u_1}} \alpha_i < \alpha_{u_1}. \quad (49)$$

Observe that

$$\sum_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}} q_j(t) = A_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(0, t) - B_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(0, t)$$

$$\geq A_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(0, t) - C_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(0, t)$$

$$\geq \left(\sum_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}} \alpha_i - F(u_1 \cup \widetilde{\mathcal{D}}_{u_1})\right) t - \mathbb{O},$$

and taking into account (49) we conclude that

$$\lim_{t \to \infty} \sum_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}} q_j(t) = \infty.$$

Since by (47) $q_i(t)$ is finite for $i \in \widetilde{\mathcal{D}}_{u_1} \subseteq \mathcal{I}_1^s$, we conclude that $\lim_{t \to \infty} q_{u_1}(t) = \infty$.

According to Corollary 6, for any subset $\mathcal{S} \subseteq \widetilde{\mathcal{D}}_{u_1}$ it holds

$$\sum_{\mathcal{S}} \alpha_i > F\left(u_1 \cup \widetilde{\mathcal{D}}_{u_1}\right) - F\left(u_1 \cup \left(\widetilde{\mathcal{D}}_{u_1} - \mathcal{S}\right)\right).$$

Notice that since $\widetilde{\mathcal{D}}_{u_1} \subseteq \mathcal{I}_1^s$, by the definition of the queue indices $I_i(q)$ the queues in $\widetilde{\mathcal{D}}_{u_1}$ are served according to the LCQ policy when they are smaller than $H = (N + 1 - u_1)T$. Moreover, since $\lim_{t \to \infty} q_{u_1}(t) = \infty$, it holds that $I_{u_1}(q_{u_1}(t)) = H$, for t larger than or equal to some time t_0. Hence, for $t \geq t_0$ queue u_1 has priority over the queues in $\mathcal{S} \subseteq \widetilde{\mathcal{D}}_{u_1}$ whenever $\max_{i \in \mathcal{S}} \{q_i(t)\} \leq H$. Therefore, we can apply Lemma 4 with $\mathcal{L} = u_1$, $\mathcal{D} = \widetilde{\mathcal{D}}_{u_1}$ and $H = (N + 1 - u_1)T$ to conclude that there is some time $\tau_0^1 \geq t_0$ such that for $i \in \widetilde{\mathcal{D}}_{u_1}$, and for all $t \geq \tau_0^1$, it holds

$$\min_{i \in \widetilde{\mathcal{D}}_{u_1}} \{q_i(t)\} \geq (N + 1 - u_1)T - \mathbb{O}. \qquad (50)$$

From Lemma 5, e) it follows that for any subset $\mathcal{S} \subseteq \mathcal{E}_1 = \mathcal{I}_1^s - \widetilde{\mathcal{D}}_{u_1}$ ($\mathcal{E}_1 = \mathcal{I}_1^s \cup \mathcal{I}_2^s - \widetilde{\mathcal{D}}_{u_1}$ if u_1 is the only element in $\mathcal{I}_1^u$) it holds

$$\sum_{\mathcal{S}} \alpha_i \leq F\left(u_1 \cup \widetilde{\mathcal{D}}_{u_1} \cup \mathcal{S}\right) - F\left(u_1 \cup \widetilde{\mathcal{D}}_{u_1}\right).$$

Applying now Lemma 2 with $\mathcal{D} = \mathcal{E}_1$, $H = (N + 1 - u_2)T$, $\Phi = T$ and $\mathcal{L} = u_1 \cup \widetilde{\mathcal{D}}_{u_1}$, we conclude that for all $t \geq 0$,

$$\max_{i \in \mathcal{E}_1} \{q_i(t)\} \leq (N + 1 - u_2)T + \mathbb{O} \qquad (51)$$

provided that $T \geq \mathbb{O}$. Pick now T large enough so that

$$T - \mathbb{O} > \mathbb{O}. \qquad (52)$$

Then since $u_2 > u_1$, it holds

$$(N + 1 - u_1)T - \mathbb{O} > (N + 1 - u_2)T + \mathbb{O}. \qquad (53)$$

Inequalities (50), (51) and (53) and the fact that $q_{u_1}(t) \to \infty$, (i.e., $I_{u_1}(q_{u_1}(t)) = (N + 1 - u_1)T$, for $t \geq \tau_0^1$) imply that the queues in $u_1 \cup \widetilde{\mathcal{D}}_{u_1}$ have higher priority over the rest of the queues for $t \geq \tau_0^1$ and that they are nonempty. Therefore, the queues in $u_1 \cup \widetilde{\mathcal{D}}_{u_1}$ use all the available channel slots. Since

$$B_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(\tau_0^1, t) = C_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(\tau_0^1, t),$$

we conclude that

$$\lim_{t \to \infty} \frac{B_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(0, t)}{t} = \lim_{t \to \infty} \frac{C_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(0, t)}{t}.$$

Taking into account (48) and the fact that

$$\lim_{t \to \infty} \frac{C_{u_1 \cup \widetilde{\mathcal{D}}_{u_1}}(0, t)}{t} = F\left(u_1 \cup \widetilde{\mathcal{D}}_{u_1}\right),$$

we conclude that $\lim_{t \to \infty}(B_{u_1}(0, t)/t)$ exists and

$$r_{u_1} = \lim_{t \to \infty} \frac{B_{u_1}(0, t)}{t} = F\left(u_1 \cup \widehat{\mathcal{D}}_{u_1}\right) - \sum_{\widehat{\mathcal{D}}_{u_1}} \alpha_i,$$

that is, $r_{u_1} = x_{u_1}^*$.

Using similar arguments we can show the claims for the rest of the queue indices. ■

V. Conclusions

We presented a policy for scheduling packets for transmission over channels with time varying connectivity. We showed that the proposed policy maximizes the weighted sum of channel throughputs under any packet arrival rates. We adopted a burstiness-constrained channel model for the analysis. This model facilitates the analysis, while at the same time allows us to dispose of statistical assumptions.

In this paper, we restricted ourselves to the case where the channels are in one of two states, "on" and "off". A subject of further work is to generalize our approach to include the case of multi-rate channels. Another subject of ongoing work is to generalize our approach to include more general optimization functions. Furthermore, the consideration of packet delays, in addition to throughput, is a practical matter that needs to be addressed.

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